

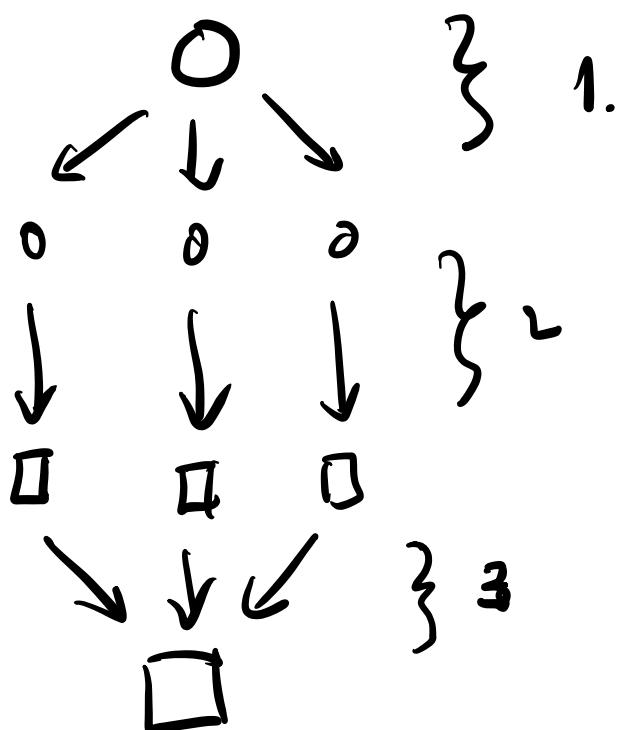
LECTURE 6 : DIVIDE & CONQUER

"Divide et impera."

- Philip II of Macedonia
382-336 BC

Divide & conquer is a strategy with 3 steps:

1. Divide a problem into subproblems
2. Turn subproblems into subolutions
3. Conquer subolutions into a solution.



```

> msort :: [Int] → [Int]
> msort [] = []
> msort [x] = [x] conquer
> msort xs = merge (msort us) (msort vs)
> where
>       (us, vs) = splitAt (n'div' 2) xs
>       n        = length xs divide

```

We need to define merge:

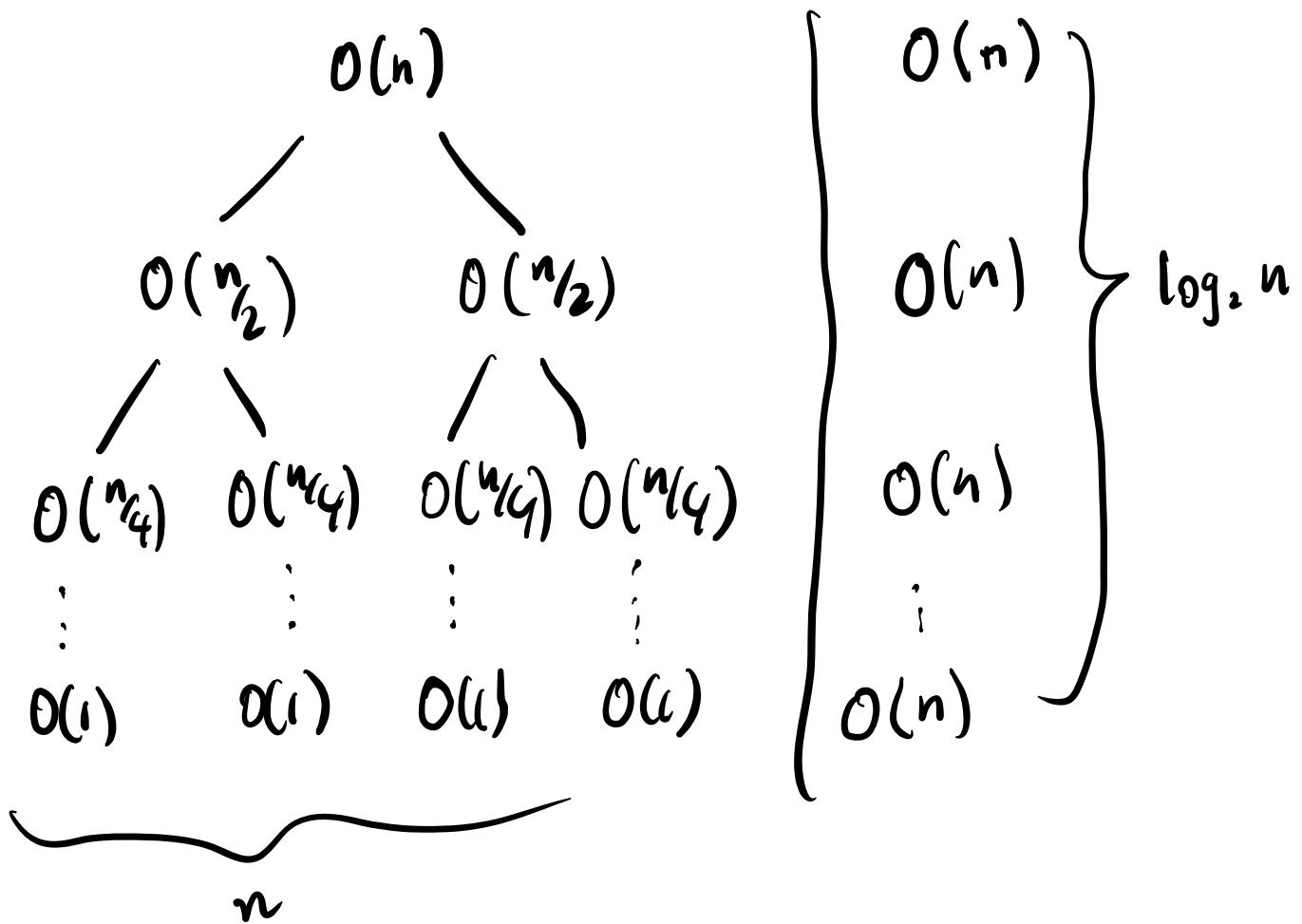
```

> merge :: [Int] → [Int] → [Int]
> merge [] ys = ys
> merge xs [] = xs
> merge xxsl (x:xs) yyrl (y:ys)
>   | x ≤ y = x : merge xs yyrl
>   | otherwise = y : merge xxsl ys

```

$$\begin{aligned}
 T(msort)(0) &= 1 \\
 T(msort)(1) &= 1 \\
 T(msort)(n) &= T(\text{length})(n) + T(\text{splitAt})\left(\frac{n}{2}\right) \\
 &\quad + 2 \cdot T(msort)(n/2)
 \end{aligned}$$

$$+ T(\text{merge})(n_2)(n_2)$$



$$T(\text{msort})(n) \in \Theta(n \cdot \log n)$$

Quicksort.

Quicksort is also a DDC algorithm

- > $\text{qsort} :: [\text{list}] \rightarrow [\text{list}]$
- > $\text{qsort} [] = []$
- > $\text{qsort} [x] = [x]$
- > $\text{qsort} (x:xs) = \text{partition}(x) \text{ } \text{conquer.}$

> $\text{qsort } us \dagger [x] \dagger \text{qsort } vs$
> where $(x: \text{qsort } vs)$

> $(us, vs) = \text{partition } (\leq x) \times s$
 \uparrow
 divide

> $\text{partition} :: (a \rightarrow \text{Bool}) \rightarrow [a] \rightarrow ([a], [a])$

> $\text{partition } p \ xs =$

> $(\text{filter } p \ xs, \text{filter } (\text{not} \cdot p) \ xs)$

Assuming the lists are split into two equal parts:

$$T(\text{qsort})(0) = 1$$

$$T(\text{qsort})(1) = 1$$

$$\begin{aligned} T(\text{qsort})(n) &= T(\text{partition})(n) \\ &+ T(\#)(\frac{n}{2}) \\ &+ 2 \cdot T(\text{qsort})(\frac{n}{2}) \end{aligned}$$

In the worst case, partition splits xs into lists of length 0 and $n-1$:

$$T(\text{qsort})(0) = 1$$

$$T(\text{qsort})(1) = 1$$

$$\begin{aligned} T(\text{qsort})(n) &= T(\text{partition})(n) \\ &+ T(\#)(n-1) \\ &+ T(\text{qsort})(0) \\ &+ T(\text{qsort})(n-1) \end{aligned}$$

$$= O(n) + T_{qsort}(n-1)$$

$$= O(n) + (O(n-1) + T_{qsort}(n-2))$$

$$= \underbrace{O(n) + \dots + O(n)}_n$$

$$= O(n^2)$$